

# Negative rank hyperoperations

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## Abstract

We construct the bi-infinite sequence of binary operations where the next element is an iteration of the previous one, such that

$$\dots, +_{r=1}, \times_{r=2}, \hat{\ }_{r=3}, \dots$$

## 1 Previous research

When constructing a sequence of binary operations that preserves the addition-multiplication-exponentiation relationship, we must choose which of these properties to construct the sequence based on.

### Hyperoperations (Goodstein's hyperoperations)

The hyperoperation sequence  $([r])_{r=0}^{\infty}$  is the sequence of binary operations  $[r] : \mathbb{N}_0 \times \mathbb{N}_0 \rightarrow \mathbb{N}_0$  defined as

$$m[r]n = \begin{cases} n + 1 & \text{if } r = 0 \\ m & \text{if } r = 1, n = 0 \\ 0 & \text{if } r = 2, n = 0 \\ 1 & \text{if } r = 3, n = 0 \\ m[r-1](m[r](n-1)) & \text{otherwise.} \end{cases}$$

This definition gives the following:

$$m[0]n = n+1, m[1]n = m+n, m[2]n = m \times n, m[3]n = m^n, m[4]n = \underbrace{m^{\dots^m}}_{n \text{ copies of } m}, \dots$$

Note that

$$m[r+1]n = \underbrace{m[r]m[r] \cdots [r]m}_{n \text{ copies of } m} \quad (\text{right-associative}) \quad (1)$$

where  $r \geq 1$ , but the iteration is one time more than, as

$$m[1]n = \underbrace{m[0]m[0] \cdots [0]m}_{n+1 \text{ copies of } m} \quad (\text{right-associative})$$

where  $r = 0$ .

## Rubtsov and Trapmann's zeration

There are some attempts to construct a new hyperoperation sequence such that equation (1) holds in general.

Let  $\star : \mathbb{N}_0 \times \mathbb{N}_2 \rightarrow \mathbb{N}_0$  be a binary operation.

$$m + n = \underbrace{m \star m \star \cdots \star m}_{n \text{ copies of } m} \quad (\text{right-associative})$$

if and only if

$$m \star n = \begin{cases} n + 2 & \text{if } n = m \\ n + 1 & \text{if } n \geq m + 2 \end{cases} \quad [1][5].$$

We consider only where  $n \geq 2$  because  $m + 1 \neq m$  (1 copy of  $m$ ). Rubtsov's zeration  $[0_R]$  defined as the following fulfills this:

$$m[0_R]n = \begin{cases} n + 2 & \text{if } n = m \\ \max(m, n) + 1 & \text{if } n \neq m. \end{cases}$$

Rubtsov's zeration is commutative  $m[0_R]n = n[0_R]m$ , distributive under addition  $l + (m[0_R]n) = (l + m)[0_R](l + n)$ , and

$$m + n = \underbrace{m[0_R]m[0_R] \cdots [0_R]m}_{n \text{ copies of } m} \quad (\text{left-associative}) \quad [6][7].$$

Also, the exact same operation has been reinvented under the name "junkasan" (which means quasi-addition)[8].

Trapmann's zeration is also a solution of the equation (1)[2]:

$$m[0_T]n = \max(m + 2, n + 1).$$

This zeration is continuous, different from Rubtsov's one.

There is the attempt to construct a bi-infinite sequence of hyperoperations such that all two adjacent elements satisfy the equation (1)[9].

## Lower hyperoperations

Another version of hyperoperation sequence that is applied left-associativity.

## Tropical semiring

The binary operation  $\min$  is commutative, associative, and distributive under addition:  $l + \min(m, n) = \min(l + m, l + n)$ . We see the finite sequence  $\min, +, \times$  is an extension of the sequence  $+, \times$  that preserves some properties of a commutative ring.

There is the attempt to construct a bi-infinite sequence of binary operations such that all two adjacent elements form a ring[10].

## Commutative hyperoperations (Bennett's hyperoperations)

Commutative hyperoperations are the sequence of binary operations defined recursively as

$$m[1]n = m + n, \quad m[r + 1]n = e^{(\ln m[r] \ln n)}.$$

This is a generalization of the following relationship between  $+$  and  $\times$  [11]:

$$m \times n = e^{(\ln m + \ln n)}.$$

The solution is  $m[r]n = \exp^{r-1}(\exp^{-(r-1)}(m) + \exp^{-(r-1)}(n))$ , and all the next element is distributive over the previous one[3]:  $l[r + 1](m[r]n) = (l[r + 1]m)[r](l[r + 1]n)$ .

Also, we can easily extend it to bi-infinite sequence. The 0th commutative hyperoperation  $m[0]n = \ln(e^m + e^n)$  is distributive under addition[4]:  $l + (m[0]n) = (l + m)[0](l + n)$ .

## Pisa hyperoperations

Pisa hyperoperations are the sequence of binary operations defined recursively as

$$m[2]n = m \times n, \quad m[r + 1]n = (m^{\frac{1}{m}})^{(m[r]n)}.$$

This is a generalization of the following relationship between  $\times$  and  $\hat{\cdot}$ :

$$m^n = (m^{\frac{1}{m}})^{(m \hat{\cdot} n)}.$$

The  $r - 1$ th operation is expressed as  $m[r - 1]n = m \ln_m(m[r]n)$ , and we can easily extend it to bi-infinite sequence[12].

## Hyperoperations we construct here

We construct the same as Eners49's one[9], the bi-infinite sequence such that  $m[1]n = m + n$  and for all  $r \in \mathbb{Z}$

$$m[r + 1]n = \underbrace{m[r]m[r] \cdots [r]m}_{n \text{ copies of } m} \quad (\text{right-associative}).$$

From here, consecutive operations are executed from the right unless otherwise stated.

Now, Rubtsov and Trapmann's zerations seem to fit for the 0th element of the bi-infinite hyperoperation sequence. However, a binary operation  $\heartsuit$  that meets the following does not exist[7]:

$$m[0_R]n, \text{ nor } m[0_T]n = \underbrace{m\heartsuit m\heartsuit \dots \heartsuit m}_{n \text{ copies of } m}.$$

*Proof.* Seeking a contradiction, suppose that

$$m[0_R]n = \underbrace{m\heartsuit m\heartsuit \dots \heartsuit m}_{n \text{ copies of } m} \quad \text{for any } m \in \mathbb{N}_0 \text{ and } n \in \mathbb{N}_2.$$

Then where  $m \geq 1$ ,

$$\begin{aligned} m[0_R](m+1) &= \underbrace{m\heartsuit m\heartsuit \dots \heartsuit m}_{m+1 \text{ copies of } m} \\ &= m\heartsuit \underbrace{(m\heartsuit m\heartsuit \dots \heartsuit m)}_{m \text{ copies of } m} \\ &= m\heartsuit(m[0_R]m) \\ &= m\heartsuit(m+2). \end{aligned} \tag{2}$$

Also,

$$\begin{aligned} m[0_R](m+2) &= \underbrace{m\heartsuit m\heartsuit \dots \heartsuit m}_{m+2 \text{ copies of } m} \\ &= m\heartsuit \underbrace{(m\heartsuit m\heartsuit \dots \heartsuit m)}_{m+1 \text{ copies of } m} \\ &= m\heartsuit(m[0_R](m+1)) \\ &= m\heartsuit(m+2). \end{aligned} \tag{3}$$

From the equation (2) and (3),

$$m+2 = m[0_R](m+1) = m\heartsuit(m+2) = m[0_R](m+2) = m+3,$$

i.e.  $2 = 3$ . □

We can use the same way to prove that about  $[0_T]$ . It is a tough question whether there exists such an operation, of which the iteration of its iteration, is equal to addition, while to find what operation repeated  $n$ -times results in addition is much easier.

## 2 The general solution of first order recurrence relations

We now want to know, for given a binary operation  $\star$ , whether an operation  $\heartsuit$  such that

$$\underbrace{m\heartsuit m\heartsuit \dots \heartsuit m}_{n \text{ copies of } m} = m\star n \quad \text{exists.}$$

Then, if  $\heartsuit$  exists, the sequence

$$m, m \star 2, m \star 3, m \star 4, \dots \quad (4)$$

is equals to

$$m, m \heartsuit m, m \heartsuit (m \heartsuit m), m \heartsuit (m \heartsuit (m \heartsuit m)), m \heartsuit (m \heartsuit (m \heartsuit (m \heartsuit m))), \dots$$

If we set  $a_1 = m, f(n) = m \heartsuit n$ , then this sequence turns into

$$a_1, f(a_1), f(f(a_1)), f(f(f(a_1))), \dots$$

That is, whether  $\heartsuit$  exists is equivalent to whether the sequence (4) is defined by a first order recurrence relation (in other words, the sequence is a dynamical system).

From here, Let  $e$  be a integer,  $\mathbb{N}_e$  be integers not less than  $e$ ,  $s$  be the successor function, and  $S$  be a set. Let  $P(a, k, k')$  be the statement  $a$  is a sequence  $\mathbb{N}_e \rightarrow S$ ,  $k, k' \in \mathbb{N}_e$ , and

$$k < k' \wedge a_k = a_{k'}.$$

$Q(a, p, k)$  be the statement  $a$  is a sequence  $\mathbb{N}_e \rightarrow S$ ,  $p \in \mathbb{N}_1$ ,  $k \in \mathbb{N}_e$ ,

$$\forall m \forall n; a_{m+k} = a_{pn+m+k} \quad (m, n \in \mathbb{N}_0).$$

**Lemma 2.1.** *Let  $f$  be a map defined on  $S$ . Suppose a sequence  $a : \mathbb{N}_e \rightarrow S$  is defined by the first order recurrence relation*

$$a_{n+1} = f(a_n).$$

*Then*

$$\forall m \in \mathbb{N}_0 \forall k \in \mathbb{N}_e; a_{m+k} = f^m(a_k).$$

*Proof.* For all  $k \in \mathbb{N}_e$ , the following holds:

Proof is by induction on  $m$ , starting at  $m = 0$ .

For  $m = 0$ ,

$$f^0(a_n) = a_n = a_{0+n}$$

and so the result holds true for  $m = 0$ .

Suppose the result holds true for some  $m$ , i.e. that  $f^m(a_n) = a_{m+n}$ . Then we have

$$f^{m+1}(a_n) = f(f^m(a_n)) = f(a_{m+n}) = a_{m+n+1} = a_{(m+1)+n}$$

and so the result holds true for  $m+1$  and our proof by induction is complete.  $\square$

**Lemma 2.2.**

*$a$  is defined by a first order recurrence relation*

$$\begin{aligned} & \Downarrow \\ & (\forall k \forall k'; P(a, k, k') \implies Q(a, k' - k, k)). \end{aligned}$$

*Proof.* By the definition of  $a$ , there are some map  $f$  defined on  $S$  such that

$$\forall a_n; a_{n+1} = f(a_n).$$

Let  $k, k' \in \mathbb{N}_e$  be

$$k < k' \wedge a_k = a_{k'}.$$

Then we see  $k' - k \in \mathbb{N}_0$ . By lemma 2.1,

$$\forall k \in \mathbb{N}_e \forall k' \in \mathbb{N}_e; a_k = a_{k'} = a_{k'-k+k} = f^{k'-k}(a_k) \quad (5)$$

For all  $m, n \in \mathbb{N}_0$ ,

$$\begin{aligned} a_{(k'-k)n+m+k} &= f^{(k'-k)n+m}(a_k) && \text{By lemma 2.1} \\ &= f^{(k'-k)n+m}(f^{k'-k}(a_k)) && \text{By (5)} \\ &= f^{(k'-k)n+m+k'-k}(a_k) && \\ &= f^{(k'-k)(n+1)+m}(a_k) && \\ &= a_{(k'-k)(n+1)+m+k} && \end{aligned} \quad (6)$$

Also,

$$a_{(k'-k)0+m+k} = a_{m+k} \quad (7)$$

From (6) and (7),

$$\forall m \forall n; a_{m+k} = a_{(k'-k)n+m+k}.$$

□

**Lemma 2.3.**

$$Q(a, p_1, k_1) \wedge Q(a, p_2, k_2) \wedge \dots \implies Q(a, \gcd(p_1, p_2, \dots), \min(k_1, k_2, \dots)).$$

That is

$$Q(a, \gcd(\pi_1[T]), \min(\pi_2[T])),$$

where  $T = \{(p, k) | Q(a, p, k)\}$ ,  $\pi_1(t)$  and  $\pi_2(t)$  are the first and second coordinate of a pair  $t$ .

*Proof.* Suppose  $p, q \in \mathbb{N}_1$ ,  $g$  is the greatest common divisor of  $p$  and  $q$ ,  $k, l \in \mathbb{N}_e$  are  $k \leq l$ ,  $Q(a, p, k)$ , and  $Q(a, q, l)$ .

Then we see that for all  $m, n \in \mathbb{N}_0$ ,

$$\begin{aligned} a_{pn+m+l} &= a_{pn+m+l-k+k} \\ &= a_{pn+(m+l-k)+k} \\ &= a_{(m+l-k)+k} && \text{by } Q(a, q, l), m+l-k \in \mathbb{N}_0 \\ &= a_{m+l}. \end{aligned} \quad (8)$$

We now check that the linear Diophantine equation

$$px - qy = r$$

has integer solutions  $(x, y)$  if and only if  $r$  is a multiple of the greatest common divisor of  $p$  and  $q$ . Then, since  $-pq < 0$ , it has a positive integer solution.

Moreover, we see  $p0 - q0 = g0$ . Hence for every  $n \in \mathbb{N}_0$ , there are  $n_1, n_2 \in \mathbb{N}_0$  such that

$$pn_1 = qn_2 + gn. \quad (9)$$

Thus for every  $m, n \in \mathbb{N}_0$ , there are  $n_1, n_2 \in \mathbb{N}_0$  such that

$$\begin{aligned} a_{gn+m+l} &= a_{(gn+m)+l} \\ &= a_{qn_2+(gn+m)+l} \quad \text{by } Q(a, q, l), gn+m \in \mathbb{N}_0 \\ &= a_{qn_2+gn+m+l} \\ &= a_{pn_1+m+l} \quad \text{by (9)} \\ &= a_{m+l} \quad \text{by (8)}. \end{aligned} \quad (10)$$

By the Archimedean property, there is  $n_3 \in \mathbb{N}_0$  such that

$$pn_3 + k \geq l \quad (11)$$

Therefore, for every  $m, n \in \mathbb{N}_0$ , there is  $n_3 \in \mathbb{N}_0$  such that

$$\begin{aligned} a_{gn+m+k} &= a_{pn_3+gn+m+k} \quad \text{by } Q(a, p, k), gn+m \in \mathbb{N}_0 \\ &= a_{pn_3+gn+m+k-l+l} \\ &= a_{gn+(pn_3+m+k-l)+l} \\ &= a_{(pn_3+m+k-l)+l} \quad \text{by (10), } pn_3+m+k-l \in \mathbb{N}_0 \text{ (by(11))} \\ &= a_{pn_3+m+k} \\ &= a_{m+k} \quad \text{by } Q(a, p, k). \end{aligned}$$

That is,

$$Q(p_1, k_1) \wedge Q(p_2, k_2) \implies Q(\gcd(p_1, p_2), \min(k_1, k_2)).$$

Therefore

$$Q(a, \gcd(\pi_1[T]), \min(\pi_2[T])),$$

where  $T = \{(p, k) | Q(a, p, k)\}$ . □

**Lemma 2.4.**

$$\begin{aligned} (\forall k \forall k'; P(a, k, k') \implies Q(a, k' - k, k)) \\ \Downarrow \\ a \text{ is injective} \vee Q(a, \gcd(\pi_1[T]), \min(\pi_2[T])), \end{aligned}$$

where  $T = \{(k' - k, k) | P(a, k, k')\}$ .

*Proof.* Where  $a$  is injective: The proposition holds.

Where  $a$  is not injective: We see  $T \neq \emptyset$ .

If  $Q(a, k' - k, k)$  holds, i.e.

$$\forall m \forall n; a_{m+k} = a_{(k'-k)n+m+k} \quad (m, n \in \mathbb{N}_0),$$

then substituting  $(m, n) = (0, 1)$ ,

$$a_{0+k} = a_{(k'-k)*1+0+k}, \text{ i.e. } a_k = a_{k'}.$$

Also we see  $k' - k \in \mathbb{N}_1$ , therefore  $k < k'$ .  $P(a, k, k')$  holds. Hence  $\forall k \forall k'; Q(a, k' - k, k) \implies P(a, k, k')$  is true.

From the hypothesis,  $\forall k \forall k'; P(a, k, k') \implies Q(a, k' - k, k)$ .

Now we have

$$\forall k \forall k'; P(a, k, k') \iff Q(a, k' - k, k).$$

Thus

$$T = \{(k' - k, k) \mid P(a, k, k')\} = \{(k' - k, k) \mid Q(a, k' - k, k)\}.$$

By lemma 2.3,

$$Q(a, \gcd(\pi_1[T]), \min(\pi_2[T])).$$

□

**Definition 2.5.** For a sequence  $a : \mathbb{N}_e \rightarrow S$ , define a set  $U(a) \subset \mathbb{N}_e$  as

$$U(a) := \{n \mid \forall m; m \leq n \implies a|_{\{|l| \leq m\}} \text{ is injective}\}.$$

That is, where  $n_1 \in \mathbb{N}_e$  is the maximum integer such that the restriction  $a|_{\{|l| \leq n_1\}}$  is injective,

$$U(a) = \begin{cases} \mathbb{N}_e & \text{if } a \text{ is injective} \\ \{n \mid n \leq n_1\} & \text{if } a \text{ is not injective.} \end{cases}$$

In addition, define  $a'$  as a map  $a' : U(a) \rightarrow a[U(a)]; n \mapsto a_n$ .

For example, for a constant sequence  $a_n = C$ ,  $a'$  becomes a map  $a' : \{e\} \rightarrow a[\{e\}]; e \mapsto a_e$  and is bijective. For all sequences  $a$ ,  $U(a)$  exists and  $a'$  is bijective.

**Lemma 2.6.** Suppose a sequence  $a : \mathbb{N}_e \rightarrow S$  is

$$a \text{ is injective} \vee Q(a, g_T, m_T).$$

Then

$$a[U(a)] = a[\mathbb{N}_e],$$

and

$$U(a) = \begin{cases} \mathbb{N}_e & \text{if } a \text{ is injective} \\ \{n \mid n \leq g_T - 1 + m_T\} & \text{if } a \text{ is not injective,} \end{cases}$$

where  $T = \{(k' - k, k) \mid P(a, k, k')\}$ ,  $g_T = \gcd(\pi_1[T])$ , and  $m_T = \min(\pi_2[T])$ .

*Proof.* Where  $a$  is injective: obvious.

Where  $a$  is not injective: We see  $Q(a, g_T, m_T)$  holds. Also  $T \neq \emptyset$ .

Let  $V$  be a set  $\{n \mid n \in \mathbb{N}_e, n \leq g_T - 1 + m_T\}$ .

Suppose the restriction  $a|_V$  is not injective. Then:

There exists  $k_1, k'_1 \in V$  such that  $k_1 < k'_1 \wedge a_{k_1} = a_{k'_1}$ .

That is,  $P(a, k_1, k'_1)$  is true.



Therefore  $(k'_1 - k_1, k_1) \in T$ ,  $k'_1 - k_1 \in \pi_1[T]$ ,  $k_1 \in \pi_2[T]$ .

For all  $k_1$ ,  $k_1 \leq m_T - 1$  or  $k_1 \geq m_T$ .

Where  $k_1 \leq m_T - 1$ :

We have  $k_1 \in \pi_2[T]$ , contradicting  $m_T = \min(\pi_2[T])$ .

Where  $k_1 \geq m_T$ :

$$-k_1 \leq -m_T.$$

$$k'_1 \leq g_T - 1 + m_T \text{ by } k'_1 \in V.$$

$$\text{Therefore } k'_1 - k_1 \leq g_T - 1 + m_T - m_T = g_T - 1.$$

We see  $k'_1 - k_1 \in \pi_1[T]$ , contradicting  $g_T = \gcd(\pi_1[T])$ .

Thus the assumption  $a|_V$  is not injective leads to a contradiction.

Hence  $a|_V$  is injective.

Let  $V'$  be a set  $\{n | n \in \mathbb{N}_e, n \leq g_T + m_T\}$ . Then we have  $m_T, g_T + m_T \in V'$ .  
By  $Q(a, g_T, m_T)$ ,  $a_{m_T} = a_{g_T + m_T}$ . Therefore

$$m_T \neq g_T + m_T \wedge a_{m_T} = a_{g_T + m_T}.$$

The restriction  $a|_{V'}$  is not injective.

By definition 2.5,

$$U(a) = V = \{n | n \in \mathbb{N}_e, n \leq g_T - 1 + m_T\}.$$

For all  $a_n \in a[\mathbb{N}_e]$ ,  $n \leq m_T - 1$  or  $n \geq m_T$ .

Where  $n \leq m_T - 1$ :

By  $n \in U(a)$ ,  $a_n \in a[U(a)]$ .

Where  $n \geq m_T$ :

There exists  $n_1, n_2 \in \mathbb{N}_0$  such that

$$n = g_T n_1 + n_2 + m_T \wedge n_2 \leq g_T - 1.$$

By  $n_2 \leq g_T - 1$ ,  $n_2 + m_T \leq g_T - 1 + m_T$ .

$$n_2 + m_T \in U(a).$$

From  $Q(a, g_T, m_T)$ ,  $a_n = a_{g_T n_1 + n_2 + m_T} = a_{n_2 + m_T} \in a[U(a)]$ .

Thus for all  $a_n \in a[\mathbb{N}_e]$ ,  $a_n \in a[U(a)]$ .

$a[\mathbb{N}_e] \subset a[U(a)]$ . Additionally, By  $U(a) \subset \mathbb{N}_e$ ,  $a[U(a)] \subset a[\mathbb{N}_e]$ . Hence  $a[U(a)] = a[\mathbb{N}_e]$ .  $\square$

**Lemma 2.7.**

$a$  is injective  $\vee Q(a, g_T, m_T) \implies a$  is defined by a first recurrence relation,

where  $T = \{(k' - k, k) | P(a, k, k')\}$ ,  $g_T = \gcd(\pi_1[T])$ , and  $m_T = \min(\pi_2[T])$ .

*Proof.* By lemma 2.6,  $a[U(a)] = a[\mathbb{N}_e]$ .

For all  $a_n \in a[\mathbb{N}_e]$ , There exists  $n_1 \in U(a)$  such that  $a_n = a_{n_1}$ .

$$\begin{aligned} a_{n+1} &= a_{n_1+1} = (a \circ s)(n_1) = (a \circ s \circ a'^{-1} \circ a')(n_1) \\ &= (a \circ s \circ a'^{-1})(a'(n_1)) = (a \circ s \circ a'^{-1})(a_{n_1}) = (a \circ s \circ a'^{-1})(a_n). \end{aligned}$$

The sequence  $a$  is defined by the recurrence relation

$$a_{n+1} = (a \circ s \circ a'^{-1})(a_n).$$

□

By lemma 2.2, 2.4, and 2.7,

$$\begin{aligned} &a \text{ is defined by a first recurrence relation} \\ &\quad \Downarrow \\ &\forall k \forall k'; P(a, k, k') \implies Q(a, k' - k, k) \\ &\quad \Downarrow \\ &a \text{ is injective} \vee Q(a, g_T, m_T) \\ &\quad \Downarrow \\ &a \text{ is defined by a first recurrence relation} \end{aligned}$$

These three propositions are equivalent.

**Lemma 2.8.** *Suppose a sequence  $a : \mathbb{N}_e \rightarrow S$  is defined by a first order recurrence relation. Let  $f$  be a map defined on  $S$ . Then*

$$\forall a_n \in a[\mathbb{N}_e]; a_{n+1} = f(a_n) \iff \forall x \in a[\mathbb{N}_e]; f(x) = a \circ s \circ a'^{-1}(x).$$

*Proof. Proof of ( $\implies$ ):* By lemma 2.6,  $a[U(a)] = a[\mathbb{N}_e]$ . For all  $x \in a[\mathbb{N}_e]$ ,  $x \in a[U(a)]$ . There exists  $n_1 \in U(a)$  such that  $x = a_{n_1}$ .

$$\begin{aligned} f(x) &= f(a_{n_1}) = a_{n_1+1} = (a \circ s)(n_1) = (a \circ s \circ a'^{-1} \circ a')(n_1) \\ &= (a \circ s \circ a'^{-1})(a'(n_1)) = (a \circ s \circ a'^{-1})(a_{n_1}) = (a \circ s \circ a'^{-1})(x). \end{aligned}$$

*(Proof of  $\impliedby$ ):* By lemma 2.6,  $a[U(a)] = a[\mathbb{N}_e]$ . For all  $a_n \in a[\mathbb{N}_e]$ , There exists  $n_1 \in U(a)$  such that  $a_n = a_{n_1}$ .

$$\begin{aligned} f(a_n) &= (a \circ s \circ a'^{-1})(a_n) = (a \circ s \circ a'^{-1})(a_{n_1}) = (a \circ s \circ a'^{-1})(a'(n_1)) \\ &= (a \circ s \circ a'^{-1} \circ a')(n_1) = (a \circ s)(n_1) = (a \circ s)(n) = a_{n+1}. \end{aligned}$$

□

If we replace "a is a sequence defined by a first order recurrence relation" in lemma 2.8 with simply "a is a sequence", then the proposition becomes false. Counterexample: the sequence

$$1, 2, 1, 1, 1, 1, \dots$$

### 3 Definition of *multiplicative* and *additive*

We now decide the domain and the codomain of operations we deal with in this article. Suppose binary operations  $\star, \heartsuit : \mathbb{N}_0 \times \mathbb{N}_0 \rightarrow \mathbb{N}_0$ , which does not necessarily form a ring, satisfy

$$\underbrace{m \star m \star \cdots \star m}_{n \text{ copies of } m} = m \heartsuit n.$$

Then  $m \heartsuit 1 = m$  (1 copy of  $m$ ) for all  $m$ , no matter what operation  $\heartsuit$  is. Addition does not suit for the operation  $\heartsuit$ . Furthermore, let  $\blacktriangleright : \mathbb{N}_0 \times \mathbb{N}_0 \rightarrow \mathbb{N}_0$  satisfy

$$\underbrace{m \heartsuit m \heartsuit \cdots \heartsuit m}_{n \text{ copies of } m} = m \blacktriangleright n.$$

Then

$$1 \blacktriangleright n = \underbrace{1 \heartsuit 1 \heartsuit \cdots \heartsuit 1}_{n \text{ copies of } 1} = 1 \text{ for all } n.$$

Even multiplication does not suit for  $\blacktriangleright$ . Thus we only deal with operations whose domain is the ordered pair of integers not less than 2. Also, because the value of

$$\underbrace{m \star m \star \cdots \star m}_{n \text{ copies of } m}$$

can be undefined where  $n \geq 3$ , we decide the codomain is integers not less than 2.

**Definition 3.1.** *Multiplicative of  $\star$ :* For a binary operation  $\star : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$  and  $m, n \in \mathbb{N}_2$ , define  $m \star^{(1)} n$  as follows:

$$m \star^{(1)} n := \underbrace{m \star m \star \cdots \star m}_{n \text{ copies of } m} = \begin{cases} m \star m & \text{if } n = 2 \\ m \star (m \star^{(1)}(n-1)) & \text{if } n \geq 3. \end{cases}$$

Then  $(m, n) \mapsto m \star^{(1)} n$  is a binary operation defined on  $\mathbb{N}_2 \times \mathbb{N}_2$ . Obviously  $m \star^{(1)} n \in \mathbb{N}_2$  for all  $m, n \in \mathbb{N}_2$ . We say that the map  $\star^{(1)} : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$  is the multiplicative of  $\star$ . Uniqueness of  $\star^{(1)}$  is clear.

We call  $\star \overbrace{(1)(1) \cdots (1)}^{n \text{ copies of } (1)}$  the  $n$ th multiplicative of  $\star$ , and denote  $\star^{(n)}$ . Also,  $\star^{(0)} := \star$ .

**Definition 3.2.** *Additive of  $\star$ :* For a binary operation  $\star : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$ , a binary operation  $\star^{(-1)} : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$  is said to be an additive of  $\star$  if  $\star^{(-1)(1)} = \star$ .

We call  $\star \overbrace{(-1)(-1) \cdots (-1)}^{n \text{ copies of } (-1)}$  a  $n$ th additive of  $\star$ , and denote  $\star^{(-n)}$ .

The additive of  $\star$ ,  $\star^{(-1)}$  does not always represent a specific operation like the indefinite integral  $\int f(x)dx$ . A range of  $\star$  such that  $(\star, n) \mapsto \star^{(-n)}$  is a well-defined map appears in a later section.

For a binary operation  $\star$ , there is no relation between whether the operation  $\heartsuit$  satisfy the following exists:

$$a \star n = \underbrace{a \heartsuit a \heartsuit \cdots \heartsuit a}_{n \text{ copies of } a} \text{ for all } n \in \mathbb{N}_2,$$

and the value of  $b \star n$  ( $b \neq a$ ). For example, let  $m \star n = 2m + n - 4$ . Then the operation  $\heartsuit$  such that

$$2 \star n = \underbrace{2 \heartsuit 2 \heartsuit \cdots \heartsuit 2}_{n \text{ copies of } 2}$$

does not exist, and where  $m \geq 3$ ,

$$m \star n = \underbrace{m \heartsuit m \heartsuit \cdots \heartsuit m}_{n \text{ copies of } m}$$

for some operations  $\heartsuit$ . We will ask the existence of such operation  $\heartsuit$  for each  $m$ .

Let a sequence  $a, b : \mathbb{N}_2 \rightarrow \mathbb{N}_2$  satisfy

$$a_n = m \star n, \quad b_n = m \heartsuit n, \quad m \star n = \underbrace{m \heartsuit m \heartsuit \cdots \heartsuit m}_{n \text{ copies of } m}.$$

Then

$$a_n = \begin{cases} b_m & \text{if } n = 2 \\ b_{a_{n-1}} & \text{if } n \geq 3. \end{cases}$$

Define  $a_1 := m$ , then

$$a_n = \begin{cases} m & \text{if } n = 1 \\ b_{a_{n-1}} & \text{if } n \geq 2. \end{cases}$$

We now have  $a_n = b_{a_{n-1}}$  for all  $n \in \mathbb{N}_2$ , the original domain of  $a$ . Also, define  $a_0 := 1, b_1 := m$ , then

$$a_n = \begin{cases} 1 & \text{if } n = 0 \\ b_{a_{n-1}} & \text{if } n \geq 1. \end{cases}$$

Define  $a_{-1} := 0, b_0 := 1$ , then

$$a_n = \begin{cases} 0 & \text{if } n = -1 \\ b_{a_{n-1}} & \text{if } n \geq 0. \end{cases}$$

Define  $a_{-2} := -1, b_{-1} = 0$ , then...

Now we show that the existence of  $\heartsuit$  s.t.

$$m \star n = \underbrace{m \heartsuit m \heartsuit \cdots \heartsuit m}_{n \text{ copies of } m}$$

is equivalent to the existence of  $B$  s.t.

$$A_n = B_{A_{n-1}},$$

where bi-infinite sequence  $A, B : \mathbb{Z} \rightarrow \mathbb{Z}$  is defined as

$$A_n = \begin{cases} n+1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m \star n & \text{if } n \geq 2, \end{cases} \quad B_n = \begin{cases} n+1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m \heartsuit n & \text{if } n \geq 2. \end{cases}$$

From here, let  $P(A, e)$  be the statement  $A$  is a bi-infinite sequence  $\mathbb{Z} \rightarrow \mathbb{Z}$ ,  $A_n = n+1$  where  $n \leq e$ , and  $A_n \geq e+2$  where  $n \geq e+1$ .

**Definition 3.3.** *Multiplicative of  $A$ : For  $A$  s.t.  $P(A, 0)$ , the bi-infinite sequence  $A^{(1)} : \mathbb{Z} \rightarrow \mathbb{Z}$  is said to be the multiplicative of  $A$  if*

$$\forall n \in \mathbb{Z}; A_n^{(1)} = A_{A_{n-1}^{(1)}} \wedge A_0^{(1)} = A_0.$$

*It is structurally the same thing whatever integer the 0 is rewritten as.*

Generally  $P(A^{(1)}, 0)$  holds true.

*Proof.* Proof is by induction on  $n$ . For  $n = 0$ :

$$A_0^{(1)} = A_0 = 0 + 1 \text{ by definition 3.3.}$$

Suppose that  $A_n^{(1)} = n+1$  for some  $n \leq 0$ . Then obviously  $A_{n-1}^{(1)} \leq 0$  because  $n-1 \leq -1$ .

$$\begin{aligned} n+1 &= A_n^{(1)} && \text{by the hypothesis} \\ &= A_{A_{n-1}^{(1)}} && \text{by definition 3.3} \\ &= A_{n-1}^{(1)} + 1 && \because A_{n-1}^{(1)} \leq 0, \end{aligned}$$

i.e.  $A_{n-1}^{(1)} = (n-1) + 1$ . By mathematical induction  $A_n^{(1)} = n+1$  for any  $n \leq 0$ .

And for  $n = 1$ :

$$A_1^{(1)} = A_{A_0^{(1)}} = A_1 \geq 2 \text{ by definition 3.3.}$$

Suppose that  $A_n^{(1)} \geq 2$  for some  $n \geq 1$ . Then

$$A_{n+1}^{(1)} = A_{A_n^{(1)}} \geq 2 \text{ by definition 3.3.}$$

By mathematical induction  $A_n^{(1)} \geq 2$  for any  $n \geq 1$ . □

Because  $P(A^{(1)}, 0)$  holds true,  $A^{(1)}$  is unique. A map  $(A, n) \mapsto A^{(n)}$  is well-defined.

**Definition 3.4.** *Additive of A:* For  $A$  s.t.  $P(A, 0)$ , the bi-infinite sequence  $A^{(-1)} : \mathbb{Z} \rightarrow \mathbb{Z}$  is said to be an additive of  $A$  if

$$P(A^{(-1)}, 0) \wedge (\forall n \in \mathbb{Z}; A_n = A_{A_{n-1}}^{(-1)}) \wedge A_0^{(-1)} = A_0,$$

that is

$$P(A^{(-1)}, 0) \wedge (A^{(-1)})^{(1)} = A.$$

We included the condition that  $P(A^{(-1)}, 0)$  holds true in the definition, because there are  $A, B$  such that

$$P(A, 0) \wedge \neg P(B, 0) \wedge (\forall n \in \mathbb{Z}; A_n = B_{A_{n-1}}) \wedge B_0 = A_0,$$

e.g.

$$A_n = \begin{cases} n+1 & \text{if } n \leq 0 \\ 5 & \text{if } n = 1 \\ n+5 & \text{if } n \geq 2, \end{cases} \quad B_n = \begin{cases} n+1 & \text{if } n \leq 0 \\ 5 & \text{if } n = 1 \\ -1 & \text{if } 2 \leq n \leq 4 \\ 7 & \text{if } n = 5 \\ -1 & \text{if } n = 6 \\ n+1 & \text{if } n \geq 7. \end{cases}$$

Similar to the binary operation's ones, we call  $A^{\overbrace{(1)(1)\dots(1)}^{n \text{ copies of } (1)}}$  the  $n$ th *multiplicative* of  $A$ , denote  $A^{(n)}$ , and vice versa. Also, we define  $A^{(0)} := A$ .

## 4 Properties of *multiplicative* and *additive*

**Lemma 4.1.** *Let  $A$  be the bi-infinite sequence s.t.  $P(A, 0)$ . If  $A$  is bijective,  $A^{(-1)}$  exists and is uniquely determined as*

$$A^{(-1)} = A \circ s \circ A^{-1}.$$

Moreover,  $A^{(-1)}$  is bijective.

*Proof. (Existence):* A sequence  $a : \mathbb{N}_1 \rightarrow \mathbb{N}_2; n \mapsto A_n$  is bijective because of  $P(A, 0)$  and that  $A$  is bijective. By lemma 2.7,  $a$  is defined by a first order recurrence relation, so there exists a map  $f : \mathbb{N}_2 \rightarrow \mathbb{N}_2$  where

$$a_{n+1} = f(a_n).$$

Define a bi-infinite sequence  $B : \mathbb{Z} \rightarrow \mathbb{Z}$  as

$$B_n := \begin{cases} n+1 & \text{if } n \leq 0 \\ a_1 & \text{if } n = 1 \\ f(n) & \text{if } n \geq 2. \end{cases}$$

Then clearly  $P(B, 0)$  holds true.

Where  $n \leq 0$ :

$$A_n = n + 1 = B_n = B_{A_{n-1}}.$$

Where  $n = 1$ :

$$A_1 = a_1 = B_1 = B_{A_0}.$$

Where  $n \geq 2$ :

$$A_n = a_n = f(a_{n-1}) = B_{a_{n-1}} = B_{A_{n-1}} \quad \because a_{n-1} \geq 2.$$

Thus  $B$  is one of *additive* of  $A$ ,  $A^{(-1)}$  by definition 3.4.

**(Uniqueness):** Let  $C : \mathbb{Z} \rightarrow \mathbb{Z}$  be an *additive* of  $A$ .

Where  $n \leq 0$ :

$A$  is bijective, and  $A(n) = n + 1$  if  $n \leq 0$ ;

therefore  $A^{-1}(n) = n - 1$  if  $n \leq 1$ .

Thus  $C(n) = n + 1 = n + 1 + 1 - 1 = (A \circ s \circ A^{-1})(n) \quad \because s(A^{-1}(n)) \leq 0$ .

Where  $n = 1$ :

$$C(1) = C(A(0)) = A(1) = (A \circ s)(0) = (A \circ s \circ A^{-1})(1).$$

Where  $n \geq 2$ :

Let  $a$  be a sequence  $\mathbb{N}_1 \rightarrow \mathbb{N}_2; n \mapsto A(n)$ .

From  $C$  is an *additive* of  $A$ ,  $A(n + 1) = C(A(n))$  for all  $n \in \mathbb{N}_1 \subset \mathbb{Z}$ .

That is  $a(n + 1) = C(a(n))$  for all  $n \in \mathbb{N}_1$ .

By lemma 2.8,  $C(n) = (a \circ s \circ a'^{-1})(n)$  for all  $n \in a[\mathbb{N}_1] = \mathbb{N}_2$ .

A sequence  $a : \mathbb{N}_1 \rightarrow \mathbb{N}_2; n \mapsto A_n$  is bijective from the assumption.

By definition 2.5,  $U(a) = \mathbb{N}_2$  and  $a' = a$ .

$A(A^{-1}(n)) = n = a(a^{-1}(n)) = A(a^{-1}(n))$  for all  $n \in \mathbb{N}_2$ .

In addition,  $A$  is bijective.

Therefore  $A^{-1}(n) = a^{-1}(n)$  for all  $n \in \mathbb{N}_2$ .

$C(n) = (a \circ s \circ a'^{-1})(n) = (a \circ s \circ a^{-1})(n) = (A \circ s \circ A^{-1})(n)$  for all  $n \in \mathbb{N}_2$ .

Thus  $C$ , the *additive* of  $A$  uniquely determined as  $C = A \circ s \circ A^{-1}$ . In other words  $A^{(-1)}$  is well-defined where  $A$  is bijective.

**(Bijectiveness):** Since  $A$ ,  $A^{-1}$ , and the successor function  $s : \mathbb{Z} \rightarrow \mathbb{Z}$  is bijective, the composite function  $A \circ s \circ A^{-1} = A^{(-1)}$  is bijective.

Because of this,  $(A, n) \mapsto A^{(-n)}$  is a well-defined map.  $\square$

By the way, an *additive* of  $A$  does not exist where  $A$  is surjective and not injective.

*Proof.* Seeking a contradiction, suppose  $A^{(-1)}$  exists. Then:

Let  $a$  be a sequence  $\mathbb{N}_1 \rightarrow \mathbb{N}_2; n \mapsto A(n)$ .

$a$  is defined by a first order recurrence that  $a(n + 1) = A^{(-1)}(a(n))$ .

By lemma 2.6,  $a[U(a)] = a[\mathbb{N}_1]$ .

Because of  $P(A, 0)$ ,  $a : \mathbb{N}_1 \rightarrow \mathbb{N}_2$  is surjective and not injective.

Since  $a$  is not injective, by definition 2.5  $U(a)$  is finite.

Thus  $a[U(a)]$  is finite.

Meanwhile, since  $a$  is surjective,  $a[\mathbb{N}_1] = \mathbb{N}_2$  is infinite.

We now arrive at a contradiction. □

**Lemma 4.2.** *Suppose a binary operator  $\star : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$ , a bi-infinite sequence  $A : \mathbb{Z} \rightarrow \mathbb{Z}$ , and an integer  $m \geq 2$  satisfy that*

$$\forall n; A(n) = \begin{cases} n + 1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m \star n & \text{if } n \geq 2. \end{cases}$$

Then

$$\forall n; A^{(1)}(n) = \begin{cases} n + 1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m \star^{(1)} n & \text{if } n \geq 2. \end{cases}$$

*Proof.* Because  $m \geq 2$  and  $\forall n; m \star n \geq 2$ ,  $P(A, 0)$  holds true. Thus  $A^{(1)} : \mathbb{Z} \rightarrow \mathbb{Z}$  is defined.

Where  $n \leq 0$ :

$A^{(1)}(n) = n + 1$  because  $P(A^{(1)}, 0)$  holds true.

Where  $n = 1$ :

$A^{(1)}(1) = A(A^{(1)}(0)) = A(1) = m$ .

Where  $n \geq 2$ :

$m \star^{(1)} n \geq 2$ .

$A^{(1)}(2) = A(A^{(1)}(1)) = A(m) = m \star m = m \star^{(1)} 2$ .

Suppose  $A^{(1)}(n) = m \star^{(1)} n$  for some  $n$ . Then:

$$A^{(1)}(n + 1) = A(A^{(1)}(n)) = A(m \star^{(1)} n) = m \star (m \star^{(1)} n) = m \star^{(1)} (n + 1).$$

Therefore, by mathematical induction  $\forall n; A^{(1)}(n) = m \star^{(1)} n$ . □



**Lemma 4.3.** Suppose a binary operator  $\star : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$ , a bi-infinite sequence  $A : \mathbb{Z} \rightarrow \mathbb{Z}$ , and an integer  $m \geq 2$  satisfy that

$$\forall n \in \mathbb{Z}; A(n) = \begin{cases} n+1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m\star n & \text{if } n \geq 2. \end{cases}$$

Let  $A$  be bijective. Then a binary operator  $\heartsuit : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$  such that

$$\forall n \in \mathbb{N}_2; m\star n = \begin{cases} m\heartsuit m & \text{if } n = 2 \\ m\heartsuit(m\star(n-1)) & \text{if } n \geq 3 \end{cases}$$

exists and uniquely determined as

$$\forall n \in \mathbb{N}_2; m\heartsuit n = (A \circ s \circ A^{-1})(n).$$

That is,

$$\forall n \in \mathbb{Z}; A^{(-1)}(n) = \begin{cases} n+1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m\heartsuit n & \text{if } n \geq 2. \end{cases}$$

*Proof. (Existence):* By lemma 4.1,  $A^{(-1)}$  exists. Let a binary operation  $\heartsuit : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$  be

$$\forall n \in \mathbb{N}_2; m\heartsuit n = A^{(-1)}(n).$$

By definition 3.4,

$$m\star 2 = A(2) = A^{(-1)}(A(1)) = A^{(-1)}(m) = m\heartsuit m$$

and

$$\forall n \geq 3; m\star n = A(n) = A^{(-1)}(A(n-1)) = A^{(-1)}(m\star(n-1)) = m\heartsuit((m\star(n-1))).$$

**(Uniqueness):** Suppose a binary operation  $\heartsuit : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$  fulfills

$$\forall n \in \mathbb{N}_2; m\star n = \begin{cases} m\heartsuit m & \text{if } n = 2 \\ m\heartsuit(m\star(n-1)) & \text{if } n \geq 3. \end{cases}$$

Define a bi-infinite sequence  $C : \mathbb{Z} \rightarrow \mathbb{Z}$  as

$$\forall n \in \mathbb{Z}; C(n) := \begin{cases} n+1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m\heartsuit n & \text{if } n \geq 2. \end{cases}$$

Then clearly  $P(C, 0)$  holds true.

Where  $n \leq 0$ :

$$A(n) = n+1 = C(n) = C(A(n-1)).$$

Where  $n = 1$ :

$$A(1) = m = C(1) = C(A(0)).$$

Where  $n = 2$ :

$$A(2) = m \star 2 = m \star m = C(m) = C(A(1)).$$

Where  $n \geq 3$ :

$$m \star (n - 1) \geq 2.$$

$$A(n) = m \star n = m \star (m \star (n - 1)) = m \star (A(n - 1)) = C(A(n - 1)).$$

Thus  $C$  is the *additive* of  $A$ .

By lemma 4.1,  $C$  is uniquely determined as  $A \circ s \circ A^{-1}$ .

$$\forall n \in \mathbb{N}_2; m \star n = C(n) = (A \circ s \circ A^{-1})(n).$$

In other words  $\star^{(-1)}$  is well-defined if  $\forall m \in \mathbb{N}_2; A : \mathbb{Z} \rightarrow \mathbb{Z}$  is bijective, where

$$\forall n \in \mathbb{Z}; A(n) = \begin{cases} n + 1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m \star n & \text{if } n \geq 2. \end{cases}$$

□

## 5 The negative hyperoperations

This proposition 5.1 was proved at [1] and [5]. However we try to prove again using lemmas we have proved here.

**Proposition 5.1.** *Let  $\star$  be a binary operation  $\mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$ .*

$$\star^{(1)} = + : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2 \iff$$

$$\forall m \in \mathbb{N}_2 \forall n \in \{m\} \cup \mathbb{Z}_{\geq m+2}; m \star n = \begin{cases} m + 2 & \text{if } n = m \\ n + 1 & \text{if } n \geq m + 2 \end{cases}$$

*Proof. (Proof of  $\implies$ ):* For all  $m \in \mathbb{N}_2$ , the following holds:

Define a sequence  $a : \mathbb{N}_1 \rightarrow \mathbb{N}_2$  as

$$\forall n \in \mathbb{N}_1; a(n) := \begin{cases} m & \text{if } n = 1 \\ m + n & \text{if } n \geq 2. \end{cases}$$

From the assumption,

$$\forall n \in \mathbb{N}_2; m + n = \begin{cases} m \star m & \text{if } n = 2 \\ m \star (m + (n - 1)) & \text{if } n \geq 3. \end{cases}$$

Therefore

$$\forall n \in \mathbb{N}_1; a(n + 1) = m \star (a(n)).$$

By lemma 2.8,

$$\forall n \in a[\mathbb{N}_1]; m \star n = (a \circ s \circ a'^{-1})(n).$$

Clearly  $a[\mathbb{N}_1] = \{m\} \cup \mathbb{Z}_{\geq m+2}$ . By definition 2.5,  $U(a) = \mathbb{N}_1$ . By lemma 2.6,  $a[U(a)] = a[\mathbb{N}_1] = \{m\} \cup \mathbb{Z}_{\geq m+2}$ . The map  $a'^{-1} : \{m\} \cup \mathbb{Z}_{\geq m+2} \rightarrow \mathbb{N}_1$  is determined as

$$\forall n \in \{m\} \cup \mathbb{Z}_{\geq m+2}; a'^{-1}(n) = \begin{cases} 1 & \text{if } n = m \\ n - m & \text{if } n \geq m + 2. \end{cases}$$

Thus

$$\begin{aligned} \forall n \in \{m\} \cup \mathbb{Z}_{\geq m+2}; m \star n &= (a \circ s \circ a'^{-1})(n) = \begin{cases} (a \circ s \circ a'^{-1})(n) & \text{if } n = m \\ (a \circ s \circ a'^{-1})(n) & \text{if } n \geq m + 2 \end{cases} \\ &= \begin{cases} (a \circ s)(1) & \text{if } n = m \\ (a \circ s)(n - m) & \text{if } n \geq m + 2 \end{cases} = \begin{cases} a(2) & \text{if } n = m \\ a(n - m + 1) & \text{if } n \geq m + 2 \end{cases} \\ &= \begin{cases} m + 2 & \text{if } n = m \\ m + (n - m + 1) & \text{if } n \geq m + 2 \end{cases} (\because n - m + 1 \geq 2) = \begin{cases} m + 2 & \text{if } n = m \\ n + 1 & \text{if } n \geq m + 2. \end{cases} \end{aligned}$$

**(Proof of  $\Leftarrow$ ):** For all  $m \in \mathbb{N}_2$ , the following holds:

From the assumption,

$$m + 2 = m \star m.$$

If  $n \geq 3$ , then  $m + (n - 1) \geq m + 2$ . Therefore

$$\forall n \geq 3; m + n = (m + (n - 1)) + 1 = m \star (m + (n - 1)).$$

Thus

$$\forall m \in \mathbb{N}_2 \forall n \in \mathbb{N}_2; m + n = \begin{cases} m \star m & \text{if } n = 2 \\ m \star (m + (n - 1)) & \text{if } n \geq 3, \end{cases}$$

i.e.  $\star^{(1)} = +$ . □

**Proposition 5.2.** Define a binary operation  $\heartsuit : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$  as

$$m \heartsuit n := \begin{cases} n & \text{if } 2 \leq n \leq m - 1 \\ m + 2 & \text{if } n = m \\ m + 1 & \text{if } n = m + 1 \\ n + 1 & \text{if } n \geq m + 2. \end{cases}$$

Then for all  $m, n \in \mathbb{N}_2$ , and  $r \in \mathbb{Z}$ ,  $m \heartsuit^{(r)} n$  exists and uniquely determined.

*Proof.* For all  $m \in \mathbb{N}_2$ , the following holds:

Define a bi-infinite sequence  $A : \mathbb{Z} \rightarrow \mathbb{Z}$  as

$$\forall n \in \mathbb{Z}; A(n) := \begin{cases} n+1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m \heartsuit n & \text{if } n \geq 2. \end{cases}$$

Where  $m = 2$ :

$$\forall n \in \mathbb{Z}; A(n) = \begin{cases} n+1 & \text{if } n \leq 0 \\ 2 & \text{if } n = 1 \\ 4 & \text{if } n = 2 \\ 3 & \text{if } n = 3 \\ n+1 & \text{if } n \geq 4, \end{cases}$$

i.e.

$$\dots, -1, 0, 1, 2, 4, 3, 5, 6, 7, \dots$$

Where  $m \geq 3$ :

$$\forall n \in \mathbb{Z}; A(n) = \begin{cases} n+1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ n & \text{if } 2 \leq n \leq m-1 \\ m+2 & \text{if } n = m \\ m+1 & \text{if } n = m+1 \\ n+1 & \text{if } n \geq m+2, \end{cases}$$

i.e.

$$\dots, -1, 0, 1, m, \underbrace{2, \dots, m-1}_{\text{Where } 2 \leq n \leq m-1}, m+2, m+1, m+3, m+4, \dots$$

Obviously  $A$  is bijective for any cases.

By definition 3.3, for all  $r \geq 0$  there exists a unique  $A^{(r)}$ .

Futhermore,  $A = A^{(0)}$  is uniquely defined and bijective. Suppose there exists a unique  $A^{(r)}$  and the  $A^{(r)}$  is bijective for some  $r \leq 0$ . Then, by lemma 4.1, the same holds for  $A^{(r-1)}$ . By mathematical induction  $A^{(r)}$  is uniquely defined for all  $r \leq 0$ .

Thus  $(r, n) \mapsto A^{(r)}(n)$  is a well-defined map  $\mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$ .

We have defined  $A^{(0)}$  as

$$\forall n \in \mathbb{Z}; A^{(0)}(n) = \begin{cases} n+1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m \heartsuit^{(0)} n & \text{if } n \geq 2. \end{cases}$$

Using lemma 4.2 and 4.3, by mathematical induction

$$\forall r \in \mathbb{Z} \forall n \in \mathbb{Z}; A^{(r)}(n) = \begin{cases} n+1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m \heartsuit^{(r)} n & \text{if } n \geq 2. \end{cases}$$

Therefore

$$\forall r \in \mathbb{Z} \forall n \in \mathbb{N}_2; m \heartsuit^{(r)} n = A^{(r)}(n).$$

Since  $A^{(r)}(n)$  is uniquely defined and  $m \heartsuit^{(r)} n \in \mathbb{N}_2$  for all  $r, n \in \mathbb{Z}$ ,  $(r, n) \mapsto m \heartsuit^{(r)} n$  works as a map  $\mathbb{Z} \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$ .

The above holds for all  $m \in \mathbb{N}_2$ . The map  $\mathbb{N}_2 \times \mathbb{Z} \times \mathbb{N}_2 \rightarrow \mathbb{N}_2; (m, r, n) \mapsto m \heartsuit^{(r)} n$  is well-defined.  $\square$

By proposition 5.1,  $\heartsuit^{(1)} = +$ . For all  $m, n \in \mathbb{N}_2$ , and  $r \in \mathbb{Z}$ ;

$$m \heartsuit^{(r+1)} n = \underbrace{m \heartsuit^{(r)} m \heartsuit^{(r)} \dots \heartsuit^{(r)} m}_{n \text{ copies of } m}.$$

That is an integer extension of the hyperoperation sequence we are finding.

There exist infinitely many bi-infinite sequences of binary operations which have the same property. A binary operation  $\star : \mathbb{N}_2 \times \mathbb{N}_2 \rightarrow \mathbb{N}_2$  such that

$$m \star n = \begin{cases} m + 2 & \text{if } n = m \\ n + 1 & \text{if } n \geq m + 2 \end{cases}$$

and  $A : \mathbb{Z} \rightarrow \mathbb{Z}$  is bijective where

$$A(n) = \begin{cases} n + 1 & \text{if } n \leq 0 \\ m & \text{if } n = 1 \\ m \star n & \text{if } n \geq 2 \end{cases}$$

can be the 0th element of the bi-infinite hyperoperation sequence, e.g.

$$m \star n = \begin{cases} m - n + 1 & \text{if } 2 \leq n \leq m - 1 \\ m + 2 & \text{if } n = m \\ m + 1 & \text{if } n = m + 1 \\ n + 1 & \text{if } n \geq m + 2 \end{cases}$$

$$m \heartsuit n = \begin{cases} n + 1 & \text{if } 2 \leq n \leq m - 2 \\ 2 & \text{if } n = m - 1 \\ m + 2 & \text{if } n = m \\ m + 1 & \text{if } n = m + 1 \\ n + 1 & \text{if } n \geq m + 2 \end{cases}$$

$$m \heartsuit^{\heartsuit} n = \begin{cases} n + 1 & \text{if } 2 \leq n \leq m - 2 \\ n + 2 & \text{if } n = m - 1, m \\ 2 & \text{if } n = m + 1 \\ n + 1 & \text{if } n \geq m + 2. \end{cases}$$

## 6 Appendices

We calculate  $-1$ st to  $-5$ th elements of the bi-infinite hyperoperation sequence, where the  $0$ th element is the binary operation  $\heartsuit$  in proposition 5.2.

### Negamoniteration

$$A(n) = \begin{cases} \begin{array}{ll} n+1 \leq 1 & \text{if } m = 2 \text{ and } n \leq 0 \\ n+1 = 2 & n = 1 \\ n+2 = 4 & n = 2 \\ n = 3 & n = 3 \\ n+1 \geq 5 & n \geq 4 \end{array} \\ \begin{array}{ll} n+1 \leq 1 & \text{if } m \geq 3 \text{ and } n \leq 0 \\ m = m & n = 1 \\ n \geq 2, \leq m-1 & 2 \leq n \leq m-1 \\ n+2 = m+2 & n = m \\ n = m+1 & n = m+1 \\ n+1 \geq m+3 & n \geq m+2 \end{array} \end{cases}$$

$$A^{-1}(n) = \begin{cases} \begin{array}{ll} n-1 \leq 0 & \text{if } m = 2 \text{ and } n \leq 1 \\ n-1 = 1 & n = 2 \\ n = 3 & n = 3 \\ n-2 = 2 & n = 4 \\ n-1 \geq 4 & n \geq 5 \end{array} \\ \begin{array}{ll} n-1 \leq 0 & \text{if } m \geq 3 \text{ and } n \leq 1 \\ n \geq 2, \leq m-1 & 2 \leq n \leq m-1 \\ 1 = 1 & n = m \\ n = m+1 & n = m+1 \\ n-2 = m & n = m+2 \\ n-1 \geq m+2 & n \geq m+3 \end{array} \end{cases}$$

$$(s \circ A^{-1})(n) = \left\{ \begin{array}{ll} n \leq 1 & \text{if } m = 2 \text{ and } n \leq 1 \\ n = 2 & n = 2 \\ n + 1 = 4 & n = 3 \\ n - 1 = 3 & n = 4 \\ n \geq 5 & n \geq 5 \\ n \leq 1 & \text{if } m \geq 3 \text{ and } n \leq 1 \\ n + 1 \geq 3, \leq m & 2 \leq n \leq m - 1 \\ 2 = 2 & n = m \\ n + 1 = m + 2 & n = m + 1 \\ n - 1 = m + 1 & n = m + 2 \\ n \geq m + 3 & n \geq m + 3 \end{array} \right. = \left\{ \begin{array}{ll} n \leq 0 & \text{if } m = 2 \text{ and } n \leq 0 \\ n = 1 & n = 1 \\ n = 2 & n = 2 \\ n + 1 = 4 & n = 3 \\ n - 1 = 3 & n = 4 \\ n \geq 5 & n \geq 5 \\ n \leq 0 & \text{if } m = 3 \text{ and } n \leq 0 \\ n = 1 & n = 1 \\ n + 1 = 3 & n = 2 \\ n - 1 = 2 & n = 3 \\ n + 1 = 5 & n = 4 \\ n - 1 = 4 & n = 5 \\ n \geq 6 & n \geq 6 \\ n \leq 0 & \text{if } m \geq 4 \text{ and } n \leq 0 \\ n = 1 & n = 1 \\ n + 1 \geq 3, \leq m - 1 & 2 \leq n \leq m - 2 \\ n + 1 = m & n = m - 1 \\ 2 = 2 & n = m \\ n + 1 = m + 2 & n = m + 1 \\ n - 1 = m + 1 & n = m + 2 \\ n \geq m + 3 & n \geq m + 3 \end{array} \right.$$

$$m \heartsuit^{(-1)} n = A^{(-1)}(n) = (A \circ s \circ A^{-1})(n) = \begin{cases} & \text{if } m = 2 \text{ and} \\ n+1 \leq 1 & n \leq 0 \\ n+1 = 2 & n = 1 \\ n+2 = 4 & n = 2 \\ n+2 = 5 & n = 3 \\ n-1 = 3 & n = 4 \\ n+1 \geq 6 & n \geq 5 \\ & \text{if } m = 3 \text{ and} \\ n+1 \leq 1 & n \leq 0 \\ n+2 = 3 & n = 1 \\ n+3 = 5 & n = 2 \\ n-1 = 2 & n = 3 \\ n+2 = 6 & n = 4 \\ n-1 = 4 & n = 5 \\ n+1 \geq 7 & n \geq 6 \\ & \text{if } m \geq 4 \text{ and} \\ n+1 \leq 1 & n \leq 0 \\ m = m & n = 1 \\ n+1 \geq 3, \leq m-1 & 2 \leq n \leq m-2 \\ n+3 = m+2 & n = m-1 \\ 2 = 2 & n = m \\ n+2 = m+3 & n = m+1 \\ n-1 = m+1 & n = m+2 \\ n+1 \geq m+4 & n \geq m+3 \end{cases}$$



## Negaditeration

$$A^{(-1)-1}(n) = \left\{ \begin{array}{ll} & \text{if } m = 2 \text{ and} \\ n - 1 \leq 0 & n \leq 1 \\ n - 1 = 1 & n = 2 \\ n + 1 = 4 & n = 3 \\ n - 2 = 2 & n = 4 \\ n - 2 = 3 & n = 5 \\ n - 1 \geq 5 & n \geq 6 \\ & \text{if } m = 3 \text{ and} \\ n - 1 \leq 0 & n \leq 1 \\ n + 1 = 3 & n = 2 \\ n - 2 = 1 & n = 3 \\ n + 1 = 5 & n = 4 \\ n - 3 = 2 & n = 5 \\ n - 2 = 4 & n = 6 \\ n - 1 \geq 6 & n \geq 7 \\ & \text{if } m \geq 4 \text{ and} \\ n - 1 \leq 0 & n \leq 1 \\ m = m & n = 2 \\ n - 1 \geq 2, \leq m - 2 & 3 \leq n \leq m - 1 \\ 1 = 1 & n = m \\ n + 1 = m + 2 & n = m + 1 \\ n - 3 = m - 1 & n = m + 2 \\ n - 2 = m + 1 & n = m + 3 \\ n - 1 \geq m + 3 & n \geq m + 4 \end{array} \right.$$

$$(\text{so}A^{(-1)-1})(n) = \left\{ \begin{array}{lll}
& & \text{if } m = 2 \text{ and} \\
n & \leq 1 & n \leq 1 \\
n & = 2 & n = 2 \\
n+2 & = 5 & n = 3 \\
n-1 & = 3 & n = 4 \\
n-1 & = 4 & n = 5 \\
n & \geq 6 & n \geq 6 \\
& & \text{if } m = 3 \text{ and} \\
n & \leq 1 & n \leq 1 \\
n+2 & = 4 & n = 2 \\
n-1 & = 2 & n = 3 \\
n+2 & = 6 & n = 4 \\
n-2 & = 3 & n = 5 \\
n-1 & = 5 & n = 6 \\
n & \geq 7 & n \geq 7 \\
& & \text{if } m \geq 4 \text{ and} \\
n & \leq 1 & n \leq 1 \\
m+1 & = m+1 & n = 2 \\
n & \geq 3, \leq m-1 & 3 \leq n \leq m-1 \\
2 & = 2 & n = m \\
n+2 & = m+3 & n = m+1 \\
n-2 & = m & n = m+2 \\
n-1 & = m+2 & n = m+3 \\
n & \geq m+4 & n \geq m+4
\end{array} \right. = \left\{ \begin{array}{lll}
& & \text{if } m = 2 \text{ and} \\
n & \leq 0 & n \leq 0 \\
n & = 1 & n = 1 \\
n & = 2 & n = 2 \\
n+2 & = 5 & n = 3 \\
n-1 & = 3 & n = 4 \\
n-1 & = 4 & n = 5 \\
n & \geq 6 & n \geq 6 \\
& & \text{if } m = 3 \text{ and} \\
n & \leq 0 & n \leq 0 \\
n & = 1 & n = 1 \\
n+2 & = 4 & n = 2 \\
n-1 & = 2 & n = 3 \\
n+2 & = 6 & n = 4 \\
n-2 & = 3 & n = 5 \\
n-1 & = 5 & n = 6 \\
n & \geq 7 & n \geq 7 \\
& & \text{if } m = 4 \text{ and} \\
n & \leq 0 & n \leq 0 \\
n & = 1 & n = 1 \\
n+3 & = 5 & n = 2 \\
n & = 3 & n = 3 \\
n-2 & = 2 & n = 4 \\
n+2 & = 7 & n = 5 \\
n-2 & = 4 & n = 6 \\
n-1 & = 6 & n = 7 \\
n & \geq 8 & n \geq 8 \\
& & \text{if } m \geq 5 \text{ and} \\
n & \leq 0 & n \leq 0 \\
n & = 1 & n = 1 \\
m+1 & = m+1 & n = 2 \\
n & \geq 3, \leq m-2 & 3 \leq n \leq m-2 \\
n & = m-1 & n = m-1 \\
2 & = 2 & n = m \\
n+2 & = m+3 & n = m+1 \\
n-2 & = m & n = m+2 \\
n-1 & = m+2 & n = m+3 \\
n & \geq m+4 & n \geq m+4
\end{array} \right.$$

$$m \heartsuit^{(-2)} n = A^{(-2)}(n) = (A^{(-1)} \circ \text{so} A^{(-1)-1})(n) = \left\{ \begin{array}{ll} \text{if } m = 2 \text{ and} \\ n + 1 \leq 1 & n \leq 0 \\ n + 1 = 2 & n = 1 \\ n + 2 = 4 & n = 2 \\ n + 3 = 6 & n = 3 \\ n + 1 = 5 & n = 4 \\ n - 2 = 3 & n = 5 \\ n + 1 \geq 7 & n \geq 6 \\ \text{if } m = 3 \text{ and} \\ n + 1 \leq 1 & n \leq 0 \\ n + 2 = 3 & n = 1 \\ n + 4 = 6 & n = 2 \\ n + 2 = 5 & n = 3 \\ n + 3 = 7 & n = 4 \\ n - 3 = 2 & n = 5 \\ n - 2 = 4 & n = 6 \\ n + 1 \geq 8 & n \geq 7 \\ \text{if } m = 4 \text{ and} \\ n + 1 \leq 1 & n \leq 0 \\ n + 3 = 4 & n = 1 \\ n + 5 = 7 & n = 2 \\ n + 3 = 6 & n = 3 \\ n - 1 = 3 & n = 4 \\ n + 3 = 8 & n = 5 \\ n - 4 = 2 & n = 6 \\ n - 2 = 5 & n = 7 \\ n + 1 \geq 9 & n \geq 8 \\ \text{if } m \geq 5 \text{ and} \\ n + 1 \leq 1 & n \leq 0 \\ m = m & n = 1 \\ m + 3 = m + 3 & n = 2 \\ n + 1 \geq 4, \leq m - 1 & 3 \leq n \leq m - 2 \\ n + 3 = m + 2 & n = m - 1 \\ 3 = 3 & n = m \\ n + 3 = m + 4 & n = m + 1 \\ 2 = 2 & n = m + 2 \\ n - 2 = m + 1 & n = m + 3 \\ n + 1 \geq m + 5 & n \geq m + 4 \end{array} \right.$$

## Negatriteration

$$A^{(-2)^{-1}}(n) = \left\{ \begin{array}{ll} \begin{array}{l} n-1 \leq 0 \\ n-1 = 1 \\ n+2 = 5 \\ n-2 = 2 \\ n-1 = 4 \\ n-3 = 3 \\ n-1 \geq 6 \end{array} & \begin{array}{l} \text{if } m = 2 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n \geq 7 \end{array} \\ \begin{array}{l} n-1 \leq 0 \\ n+3 = 5 \\ n-2 = 1 \\ n+2 = 6 \\ n-2 = 3 \\ n-4 = 2 \\ n-3 = 4 \\ n-1 \geq 7 \end{array} & \begin{array}{l} \text{if } m = 3 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n \geq 8 \end{array} \\ \begin{array}{l} n-1 \leq 0 \\ n+4 = 6 \\ n+1 = 4 \\ n-3 = 1 \\ n+2 = 7 \\ n-3 = 3 \\ n-5 = 2 \\ n-3 = 5 \\ n-1 \geq 8 \end{array} & \begin{array}{l} \text{if } m = 4 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n \geq 9 \end{array} \\ \begin{array}{l} n-1 \leq 0 \\ m+2 = m+2 \\ m = m \\ n-1 \geq 3, \leq m-2 \\ 1 = 1 \\ n+2 = m+3 \\ n-3 = m-1 \\ 2 = 2 \\ n-3 = m+1 \\ n-1 \geq m+4 \end{array} & \begin{array}{l} \text{if } m \geq 5 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ 4 \leq n \leq m-1 \\ n = m \\ n = m+1 \\ n = m+2 \\ n = m+3 \\ n = m+4 \\ n \geq m+5 \end{array} \end{array} \right.$$

$$(\text{so}A^{(-2)^{-1}})(n) = \left\{ \begin{array}{ll}
\begin{array}{l}
n \leq 1 \\
n = 2 \\
n+3 = 6 \\
n-1 = 3 \\
n = 5 \\
n-2 = 4 \\
n \geq 7 \\
\\
n \leq 1 \\
n+4 = 6 \\
n-1 = 2 \\
n+3 = 7 \\
n-1 = 4 \\
n-3 = 3 \\
n-2 = 5 \\
n \geq 8 \\
\\
n \leq 1 \\
n+5 = 7 \\
n+2 = 5 \\
n-2 = 2 \\
n+3 = 8 \\
n-2 = 4 \\
n-4 = 3 \\
n-2 = 6 \\
n \geq 9 \\
\\
n \leq 1 \\
m+3 = m+3 \\
m+1 = m+1 \\
n \geq 4, \leq m-1 \\
2 = 2 \\
n+3 = m+4 \\
n-2 = m \\
3 = 3 \\
n-2 = m+2 \\
n \geq m+5
\end{array}
&
\begin{array}{l}
\text{if } m = 2 \text{ and} \\
n \leq 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n \geq 7 \\
\\
\text{if } m = 3 \text{ and} \\
n \leq 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n \geq 8 \\
\\
\text{if } m = 4 \text{ and} \\
n \leq 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n \geq 9 \\
\\
\text{if } m \geq 5 \text{ and} \\
n \leq 1 \\
n = 2 \\
n = 3 \\
4 \leq n \leq m-1 \\
n = m \\
n = m+1 \\
n = m+2 \\
n = m+3 \\
n = m+4 \\
n \geq m+5
\end{array}
\end{array} \right. = \left\{ \begin{array}{ll}
\begin{array}{l}
n \leq 0 \\
n = 1 \\
n = 2 \\
n+3 = 6 \\
n-1 = 3 \\
n = 5 \\
n-2 = 4 \\
n \geq 7 \\
\\
n \leq 0 \\
n = 1 \\
n+4 = 6 \\
n-1 = 2 \\
n+3 = 7 \\
n-1 = 4 \\
n-3 = 3 \\
n-2 = 5 \\
n \geq 8 \\
\\
n \leq 0 \\
n = 1 \\
n+5 = 7 \\
n+2 = 5 \\
n-2 = 2 \\
n+3 = 8 \\
n-2 = 4 \\
n-4 = 3 \\
n-2 = 6 \\
n \geq 9 \\
\\
n \leq 0 \\
n = 1 \\
n+6 = 8 \\
n+3 = 6 \\
n = 4 \\
n-3 = 2 \\
n+3 = 9 \\
n-2 = 5 \\
n-5 = 3 \\
n-2 = 7 \\
n \geq 10 \\
\\
n \leq 0 \\
n = 1 \\
m+3 = m+3 \\
m+1 = m+1 \\
n \geq 4, \leq m-2 \\
2 = 2 \\
n+3 = m+4 \\
n-2 = m \\
3 = 3 \\
n-2 = m+2 \\
n \geq m+5
\end{array}
&
\begin{array}{l}
\text{if } m = 2 \text{ and} \\
n \leq 0 \\
n = 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n \geq 7 \\
\\
\text{if } m = 3 \text{ and} \\
n \leq 0 \\
n = 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n \geq 8 \\
\\
\text{if } m = 4 \text{ and} \\
n \leq 0 \\
n = 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n \geq 9 \\
\\
\text{if } m = 5 \text{ and} \\
n \leq 0 \\
n = 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n = 9 \\
n \geq 10 \\
\\
\text{if } m \geq 6 \text{ and} \\
n \leq 0 \\
n = 1 \\
n = 2 \\
n = 3 \\
4 \leq n \leq m-2 \\
n = m-1 \\
n = m \\
n = m+1 \\
n = m+2 \\
n = m+3 \\
n = m+4 \\
n \geq m+5
\end{array}
\end{array} \right.$$

$$m \heartsuit^{(-3)} n = A^{(-3)}(n) = (A^{(-2)} \circ_{\text{so}} A^{(-2)-1})(n) = \left\{ \begin{array}{ll} \begin{array}{l} n+1 \leq 1 \\ n+1 = 2 \\ n+2 = 4 \\ n+4 = 7 \\ n+2 = 6 \\ n-2 = 3 \\ n-1 = 5 \\ n+1 \geq 8 \end{array} & \begin{array}{l} \text{if } m = 2 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n \geq 7 \end{array} \\ \begin{array}{l} n+1 \leq 1 \\ n+2 = 3 \\ n+2 = 4 \\ n+3 = 6 \\ n+4 = 8 \\ n+2 = 7 \\ n-1 = 5 \\ n-5 = 2 \\ n+1 \geq 9 \end{array} & \begin{array}{l} \text{if } m = 3 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n \geq 8 \end{array} \\ \begin{array}{l} n+1 \leq 1 \\ n+3 = 4 \\ n+3 = 5 \\ n+5 = 8 \\ n+3 = 7 \\ n+4 = 9 \\ n-3 = 3 \\ n-1 = 6 \\ n-6 = 2 \\ n+1 \geq 10 \end{array} & \begin{array}{l} \text{if } m = 4 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n \geq 9 \end{array} \\ \begin{array}{l} n+1 \leq 1 \\ n+4 = 5 \\ n+4 = 6 \\ n+6 = 9 \\ n+3 = 7 \\ n+3 = 8 \\ n+4 = 10 \\ n-4 = 3 \\ n-4 = 4 \\ n-7 = 2 \\ n+1 \geq 11 \end{array} & \begin{array}{l} \text{if } m = 5 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n = 9 \\ n \geq 10 \end{array} \\ \begin{array}{l} n+1 \leq 1 \\ m = m \\ m+1 = m+1 \\ m+4 = m+4 \\ n+1 \geq 5, \leq m-1 \\ n+3 = m+2 \\ n+3 = m+3 \\ n+4 = m+5 \\ 3 = 3 \\ 4 = 4 \\ 2 = 2 \\ n+1 \geq m+6 \end{array} & \begin{array}{l} \text{if } m \geq 6 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ 4 \leq n \leq m-2 \\ n = m-1 \\ n = m \\ n = m+1 \\ n = m+2 \\ n = m+3 \\ n = m+4 \\ n \geq m+5 \end{array} \end{array} \right.$$

## Negatetration

$$A^{(-3)^{-1}}(n) = \left\{ \begin{array}{ll} \begin{array}{l} n-1 \leq 0 \\ n-1 = 1 \\ n+2 = 5 \\ n-2 = 2 \\ n+1 = 6 \\ n-2 = 4 \\ n-4 = 3 \\ n-1 \geq 7 \end{array} & \begin{array}{l} \text{if } m = 2 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n \geq 8 \end{array} \\ \\ \begin{array}{l} n-1 \leq 0 \\ n+5 = 7 \\ n-2 = 1 \\ n-2 = 2 \\ n+1 = 6 \\ n-3 = 3 \\ n-2 = 5 \\ n-4 = 4 \\ n-1 \geq 8 \end{array} & \begin{array}{l} \text{if } m = 3 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n \geq 9 \end{array} \\ \\ \begin{array}{l} n-1 \leq 0 \\ n+6 = 8 \\ n+3 = 6 \\ n-3 = 1 \\ n-3 = 2 \\ n+1 = 7 \\ n-3 = 4 \\ n-5 = 3 \\ n-4 = 5 \\ n-1 \geq 9 \end{array} & \begin{array}{l} \text{if } m = 4 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n = 9 \\ n \geq 10 \end{array} \\ \\ \begin{array}{l} n-1 \leq 0 \\ n+7 = 9 \\ n+4 = 7 \\ n+4 = 8 \\ n-4 = 1 \\ n-4 = 2 \\ n-3 = 4 \\ n-3 = 5 \\ n-6 = 3 \\ n-4 = 6 \\ n-1 \geq 10 \end{array} & \begin{array}{l} \text{if } m = 5 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n = 9 \\ n = 10 \\ n \geq 11 \end{array} \\ \\ \begin{array}{l} n-1 \leq 0 \\ m+4 = m+4 \\ m+2 = m+2 \\ m+3 = m+3 \\ n-1 \geq 4, \leq m-2 \\ 1 = 1 \\ 2 = 2 \\ n-3 = m-1 \\ n-3 = m \\ 3 = 3 \\ n-4 = m+1 \\ n-1 \geq m+5 \end{array} & \begin{array}{l} \text{if } m \geq 6 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ 5 \leq n \leq m-1 \\ n = m \\ n = m+1 \\ n = m+2 \\ n = m+3 \\ n = m+4 \\ n = m+5 \\ n \geq m+6 \end{array} \end{array} \right.$$

$$(\text{so}A^{(-3)^{-1}})(n) = \left\{ \begin{array}{lll}
n & \leq 1 & \text{if } m = 2 \text{ and} \\
n & = 2 & n \leq 1 \\
n + 3 & = 6 & n = 2 \\
n - 1 & = 3 & n = 3 \\
n + 2 & = 7 & n = 4 \\
n - 1 & = 5 & n = 5 \\
n - 3 & = 4 & n = 6 \\
n & \geq 8 & n = 7 \\
& & n \geq 8 \\
n & \leq 1 & \text{if } m = 3 \text{ and} \\
n + 6 & = 8 & n \leq 1 \\
n - 1 & = 2 & n = 2 \\
n - 1 & = 3 & n = 3 \\
n + 2 & = 7 & n = 4 \\
n - 2 & = 4 & n = 5 \\
n - 1 & = 6 & n = 6 \\
n - 3 & = 5 & n = 7 \\
n & \geq 9 & n \geq 8 \\
& & \text{if } m = 4 \text{ and} \\
n & \leq 1 & n \leq 1 \\
n + 7 & = 9 & n = 2 \\
n + 4 & = 7 & n = 3 \\
n - 2 & = 2 & n = 4 \\
n - 2 & = 3 & n = 5 \\
n + 2 & = 8 & n = 6 \\
n - 2 & = 5 & n = 7 \\
n - 4 & = 4 & n = 8 \\
n - 3 & = 6 & n = 9 \\
n & \geq 10 & n \geq 10 \\
& & \text{if } m = 5 \text{ and} \\
n & \leq 1 & n \leq 1 \\
n + 8 & = 10 & n = 2 \\
n + 5 & = 8 & n = 3 \\
n + 5 & = 9 & n = 4 \\
n - 3 & = 2 & n = 5 \\
n - 3 & = 3 & n = 6 \\
n - 2 & = 5 & n = 7 \\
n - 2 & = 6 & n = 8 \\
n - 5 & = 4 & n = 9 \\
n - 3 & = 7 & n = 10 \\
n & \geq 11 & n \geq 11 \\
& & \text{if } m \geq 6 \text{ and} \\
n & \leq 1 & n \leq 1 \\
m + 5 & = m + 5 & n = 2 \\
m + 3 & = m + 3 & n = 3 \\
m + 4 & = m + 4 & n = 4 \\
n & \geq 5, \leq m - 1 & 5 \leq n \leq m - 1 \\
2 & = 2 & n = m \\
3 & = 3 & n = m + 1 \\
n - 2 & = m & n = m + 2 \\
n - 2 & = m + 1 & n = m + 3 \\
4 & = 4 & n = m + 4 \\
n - 3 & = m + 2 & n = m + 5 \\
n & \geq m + 6 & n \geq m + 6
\end{array} \right. = \left\{ \begin{array}{lll}
n & \leq 0 & \text{if } m = 2 \text{ and} \\
n & = 1 & n \leq 0 \\
n & = 2 & n = 1 \\
n + 3 & = 6 & n = 2 \\
n - 1 & = 3 & n = 3 \\
n + 2 & = 7 & n = 4 \\
n - 1 & = 5 & n = 5 \\
n - 3 & = 4 & n = 6 \\
n & \geq 8 & n \geq 7 \\
& & \text{if } m = 3 \text{ and} \\
n & \leq 0 & n \leq 0 \\
n & = 1 & n = 1 \\
n + 6 & = 8 & n = 2 \\
n - 1 & = 2 & n = 3 \\
n - 1 & = 3 & n = 4 \\
n + 2 & = 7 & n = 5 \\
n - 2 & = 4 & n = 6 \\
n - 1 & = 6 & n = 7 \\
n - 3 & = 5 & n = 8 \\
n & \geq 9 & n \geq 9 \\
& & \text{if } m = 4 \text{ and} \\
n & \leq 0 & n \leq 0 \\
n & = 1 & n = 1 \\
n + 7 & = 9 & n = 2 \\
n + 4 & = 7 & n = 3 \\
n - 2 & = 2 & n = 4 \\
n - 2 & = 3 & n = 5 \\
n + 2 & = 8 & n = 6 \\
n - 2 & = 5 & n = 7 \\
n - 4 & = 4 & n = 8 \\
n - 3 & = 6 & n = 9 \\
n & \geq 10 & n \geq 10 \\
& & \text{if } m = 5 \text{ and} \\
n & \leq 0 & n \leq 0 \\
n & = 1 & n = 1 \\
n + 8 & = 10 & n = 2 \\
n + 5 & = 8 & n = 3 \\
n + 5 & = 9 & n = 4 \\
n - 3 & = 2 & n = 5 \\
n - 3 & = 3 & n = 6 \\
n - 2 & = 5 & n = 7 \\
n - 2 & = 6 & n = 8 \\
n - 5 & = 4 & n = 9 \\
n - 3 & = 7 & n = 10 \\
n & \geq 11 & n \geq 11 \\
& & \text{if } m = 6 \text{ and} \\
n & \leq 0 & n \leq 0 \\
n & = 1 & n = 1 \\
n + 9 & = 11 & n = 2 \\
n + 6 & = 9 & n = 3 \\
n + 6 & = 10 & n = 4 \\
n & = 5 & n = 5 \\
n - 4 & = 2 & n = 6 \\
n - 4 & = 3 & n = 7 \\
n - 2 & = 6 & n = 8 \\
n - 2 & = 7 & n = 9 \\
n - 6 & = 4 & n = 10 \\
n - 3 & = 8 & n = 11 \\
n & \geq 12 & n \geq 12 \\
& & \text{if } m \geq 7 \text{ and} \\
n & \leq 0 & n \leq 0 \\
n & = 1 & n = 1 \\
m + 5 & = m + 5 & n = 2 \\
m + 3 & = m + 3 & n = 3 \\
m + 4 & = m + 4 & n = 4 \\
n & \geq 5, \leq m - 2 & 5 \leq n \leq m - 2 \\
n & = m - 1 & n = m - 1 \\
2 & = 2 & n = m \\
3 & = 3 & n = m + 1 \\
n - 2 & = m & n = m + 2 \\
n - 2 & = m + 1 & n = m + 3 \\
4 & = 4 & n = m + 4 \\
n - 3 & = m + 2 & n = m + 5 \\
n & \geq m + 6 & n \geq m + 6
\end{array} \right.$$



$$m \heartsuit^{(-4)} n = A^{(-4)}(n) = (A^{(-3)} \circ s \circ A^{(-3)-1})(n) = \left\{ \begin{array}{ll} \begin{array}{l} n+1 \leq 1 \\ n+1 = 2 \\ n+2 = 4 \\ n+2 = 5 \\ n+3 = 7 \\ n+3 = 8 \\ n-3 = 3 \\ n-1 = 6 \\ n+1 \geq 9 \end{array} & \begin{array}{l} \text{if } m = 2 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n \geq 8 \end{array} \\ \begin{array}{l} n+1 \leq 1 \\ n+2 = 3 \\ n+7 = 9 \\ n+1 = 4 \\ n+2 = 6 \\ n-3 = 2 \\ n+2 = 8 \\ n-2 = 5 \\ n-1 = 7 \\ n+1 \geq 10 \end{array} & \begin{array}{l} \text{if } m = 3 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n \geq 9 \end{array} \\ \begin{array}{l} n+1 \leq 1 \\ n+3 = 4 \\ n+8 = 10 \\ n+3 = 6 \\ n+1 = 5 \\ n+3 = 8 \\ n-4 = 2 \\ n+2 = 9 \\ n-1 = 7 \\ n-6 = 3 \\ n+1 \geq 11 \end{array} & \begin{array}{l} \text{if } m = 4 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n = 9 \\ n \geq 10 \end{array} \\ \begin{array}{l} n+1 \leq 1 \\ n+4 = 5 \\ n+9 = 11 \\ n+1 = 4 \\ n-2 = 2 \\ n+1 = 6 \\ n+3 = 9 \\ n+1 = 8 \\ n+2 = 10 \\ n-2 = 7 \\ n-7 = 3 \\ n+1 \geq 12 \end{array} & \begin{array}{l} \text{if } m = 5 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n = 9 \\ n = 10 \\ n \geq 11 \end{array} \\ \begin{array}{l} n+1 \leq 1 \\ n+5 = 6 \\ n+10 = 12 \\ n+1 = 4 \\ n-2 = 2 \\ n+3 = 8 \\ n+1 = 7 \\ n+3 = 10 \\ n+1 = 9 \\ n+2 = 11 \\ n-5 = 5 \\ n-8 = 3 \\ n+1 \geq 13 \end{array} & \begin{array}{l} \text{if } m = 6 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n = 9 \\ n = 10 \\ n = 11 \\ n \geq 12 \end{array} \\ \begin{array}{l} n+1 \leq 1 \\ m = m \\ m+6 = m+6 \\ n+1 = 4 \\ n-2 = 2 \\ n+1 \geq 6, \leq m-1 \\ n+3 = m+2 \\ n+1 = m+1 \\ n+3 = m+4 \\ n+1 = m+3 \\ n+2 = m+5 \\ 5 = 5 \\ 3 = 3 \\ n+1 \geq m+7 \end{array} & \begin{array}{l} \text{if } m \geq 7 \text{ and} \\ n \leq 0 \\ n = 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ 5 \leq n \leq m-2 \\ n = m-1 \\ n = m \\ n = m+1 \\ n = m+2 \\ n = m+3 \\ n = m+4 \\ n = m+5 \\ n \geq m+6 \end{array} \end{array} \right.$$

# Negapentation

$$A^{(-4)^{-1}}(n) = \left\{ \begin{array}{ll} \begin{array}{l} n-1 \leq 0 \\ n-1 = 1 \\ n+3 = 6 \\ n-2 = 2 \\ n-2 = 3 \\ n+1 = 7 \\ n-3 = 4 \\ n-3 = 5 \\ n-1 \geq 8 \end{array} & \begin{array}{l} \text{if } m = 2 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n \geq 9 \end{array} \\ \begin{array}{l} n-1 \leq 0 \\ n+3 = 5 \\ n-2 = 1 \\ n-1 = 3 \\ n+2 = 7 \\ n-2 = 4 \\ n+1 = 8 \\ n-2 = 6 \\ n-7 = 2 \\ n-1 \geq 9 \end{array} & \begin{array}{l} \text{if } m = 3 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n = 9 \\ n \geq 10 \end{array} \\ \begin{array}{l} n-1 \leq 0 \\ n+4 = 6 \\ n+6 = 9 \\ n-3 = 1 \\ n-1 = 4 \\ n-3 = 3 \\ n+1 = 8 \\ n-3 = 5 \\ n-2 = 7 \\ n-8 = 2 \\ n-1 \geq 10 \end{array} & \begin{array}{l} \text{if } m = 4 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n = 9 \\ n = 10 \\ n \geq 11 \end{array} \\ \begin{array}{l} n-1 \leq 0 \\ n+2 = 4 \\ n+7 = 10 \\ n-1 = 3 \\ n-4 = 1 \\ n-1 = 5 \\ n+2 = 9 \\ n-1 = 7 \\ n-3 = 6 \\ n-2 = 8 \\ n-9 = 2 \\ n-1 \geq 11 \end{array} & \begin{array}{l} \text{if } m = 5 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n = 9 \\ n = 10 \\ n = 11 \\ n \geq 12 \end{array} \\ \begin{array}{l} n-1 \leq 0 \\ n+2 = 4 \\ n+8 = 11 \\ n-1 = 3 \\ n+5 = 10 \\ n-5 = 1 \\ n-1 = 6 \\ n-3 = 5 \\ n-1 = 8 \\ n-3 = 7 \\ n-2 = 9 \\ n-10 = 2 \\ n-1 \geq 12 \end{array} & \begin{array}{l} \text{if } m = 6 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ n = 6 \\ n = 7 \\ n = 8 \\ n = 9 \\ n = 10 \\ n = 11 \\ n = 12 \\ n \geq 13 \end{array} \\ \begin{array}{l} n-1 \leq 0 \\ n+2 = 4 \\ m+5 = m+5 \\ n-1 = 3 \\ m+4 = m+4 \\ n-1 \geq 5, \leq m-2 \\ 1 = 1 \\ n-1 = m \\ n-3 = m-1 \\ n-1 = m+2 \\ n-3 = m+1 \\ n-2 = m+3 \\ 2 = 2 \\ n-1 \geq m+6 \end{array} & \begin{array}{l} \text{if } m \geq 7 \text{ and} \\ n \leq 1 \\ n = 2 \\ n = 3 \\ n = 4 \\ n = 5 \\ 6 \leq n \leq m-1 \\ n = m \\ n = m+1 \\ n = m+2 \\ n = m+3 \\ n = m+4 \\ n = m+5 \\ n = m+6 \\ n \geq m+7 \end{array} \end{array} \right.$$

$$\begin{aligned}
(s \circ A^{(-4)-1})(n) &= \left\{ \begin{array}{l}
n \leq 1 \\
n \leq 2 \\
n+4 \leq 7 \\
n-1 \leq 3 \\
n-1 \leq 4 \\
n+2 \leq 8 \\
n-2 \leq 5 \\
n-2 \leq 6 \\
n \geq 9 \\
\\
n \leq 1 \\
n+4 \leq 6 \\
n-1 \leq 2 \\
n \leq 4 \\
n+3 \leq 8 \\
n-1 \leq 5 \\
n+2 \leq 9 \\
n-1 \leq 7 \\
n-1 \leq 3 \\
n \geq 10 \\
\\
n \leq 1 \\
n+5 \leq 7 \\
n+7 \leq 10 \\
n-2 \leq 2 \\
n \leq 5 \\
n-2 \leq 4 \\
n+2 \leq 9 \\
n-2 \leq 6 \\
n-1 \leq 8 \\
n-7 \leq 3 \\
n \geq 11 \\
\\
n \leq 1 \\
n+3 \leq 5 \\
n+8 \leq 11 \\
n-3 \leq 4 \\
n-3 \leq 2 \\
n+3 \leq 6 \\
n+3 \leq 10 \\
n \leq 8 \\
n-2 \leq 7 \\
n-1 \leq 9 \\
n-8 \leq 3 \\
n \geq 12 \\
\\
n \leq 1 \\
n+3 \leq 5 \\
n+9 \leq 12 \\
n \leq 4 \\
n+6 \leq 11 \\
n-4 \leq 2 \\
n \leq 7 \\
n-2 \leq 6 \\
n \leq 9 \\
n-2 \leq 8 \\
n-1 \leq 10 \\
n-9 \leq 3 \\
n \geq 13 \\
\\
n \leq 1 \\
n+3 \leq 5 \\
m+6 \leq m+6 \\
n \leq 4 \\
m+5 \leq m+5 \\
2 \leq 6, \leq m-1 \\
n \leq 2 \\
n \leq m+1 \\
n-2 \leq m \\
n \leq m+3 \\
n-2 \leq m+2 \\
n-1 \leq m+4 \\
3 \leq m+3 \\
n \geq m+7
\end{array} \right.
\begin{array}{l}
\text{if } m = 2 \text{ and} \\
n \leq 1 \\
n \leq 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n \geq 9 \\
\text{if } m = 3 \text{ and} \\
n \leq 1 \\
n \leq 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n = 9 \\
n \geq 10 \\
\text{if } m = 4 \text{ and} \\
n \leq 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n = 9 \\
n \geq 10 \\
\text{if } m = 5 \text{ and} \\
n \leq 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n = 9 \\
n = 10 \\
n = 11 \\
n \geq 12 \\
\text{if } m = 6 \text{ and} \\
n \leq 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n = 9 \\
n = 10 \\
n = 11 \\
n = 12 \\
n \geq 13 \\
\text{if } m \geq 7 \text{ and} \\
n \leq 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
6 \leq n \leq m-1 \\
n = m \\
n = m+1 \\
n = m+2 \\
n = m+3 \\
n = m+4 \\
n = m+5 \\
n = m+6 \\
n \geq m+7
\end{array}
\right. = \left\{ \begin{array}{l}
n \leq 0 \\
n \leq 1 \\
n \leq 2 \\
n+4 \leq 7 \\
n-1 \leq 3 \\
n-1 \leq 4 \\
n+2 \leq 8 \\
n-2 \leq 5 \\
n-2 \leq 6 \\
n \geq 9 \\
\\
n \leq 0 \\
n \leq 1 \\
n+4 \leq 6 \\
n-1 \leq 2 \\
n+3 \leq 8 \\
n-1 \leq 5 \\
n+2 \leq 9 \\
n-1 \leq 7 \\
n-6 \leq 3 \\
n \geq 10 \\
\\
n \leq 0 \\
n \leq 1 \\
n+5 \leq 7 \\
n+7 \leq 10 \\
n-2 \leq 2 \\
n \leq 5 \\
n-2 \leq 4 \\
n+2 \leq 9 \\
n-2 \leq 6 \\
n-1 \leq 8 \\
n-7 \leq 3 \\
n \geq 11 \\
\\
n \leq 0 \\
n \leq 1 \\
n+3 \leq 5 \\
n+8 \leq 11 \\
n-3 \leq 4 \\
n-3 \leq 2 \\
n+3 \leq 6 \\
n+3 \leq 10 \\
n \leq 8 \\
n-2 \leq 7 \\
n-1 \leq 9 \\
n-8 \leq 3 \\
n \geq 12 \\
\\
n \leq 0 \\
n \leq 1 \\
n+3 \leq 5 \\
n+9 \leq 12 \\
n-2 \leq 8 \\
n-1 \leq 10 \\
n-9 \leq 3 \\
n \geq 13 \\
\\
n \leq 0 \\
n \leq 1 \\
n+3 \leq 5 \\
n+10 \leq 13 \\
n \leq 4 \\
n+7 \leq 12 \\
n-5 \leq 2 \\
n \leq 8 \\
n-2 \leq 7 \\
n-2 \leq 10 \\
n-1 \leq 11 \\
n-10 \leq 3 \\
n \geq 14 \\
\\
n \leq 0 \\
n \leq 1 \\
n+3 \leq 5 \\
m+6 \leq m+6 \\
n \leq 4 \\
m+5 \leq m+5 \\
n \geq 6, \leq m-2 \\
n \leq m-1 \\
2 \leq m \\
n-2 \leq m+1 \\
n \leq m \\
n-2 \leq m+3 \\
n-2 \leq m+2 \\
n-1 \leq m+4 \\
3 \leq m+3 \\
n \geq m+7
\end{array} \right.
\begin{array}{l}
\text{if } m = 2 \text{ and} \\
n \leq 0 \\
n \leq 1 \\
n \leq 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n \geq 9 \\
\text{if } m = 3 \text{ and} \\
n \leq 0 \\
n \leq 1 \\
n \leq 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n = 9 \\
n \geq 10 \\
\text{if } m = 4 \text{ and} \\
n \leq 0 \\
n \leq 1 \\
n \leq 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n = 9 \\
n \geq 10 \\
\text{if } m = 5 \text{ and} \\
n \leq 0 \\
n \leq 1 \\
n \leq 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n = 9 \\
n = 10 \\
n = 11 \\
n \geq 12 \\
\text{if } m = 6 \text{ and} \\
n \leq 0 \\
n \leq 1 \\
n \leq 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n = 9 \\
n = 10 \\
n = 11 \\
n = 12 \\
n \geq 13 \\
\text{if } m = 7 \text{ and} \\
n \leq 0 \\
n \leq 1 \\
n \leq 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
n = 6 \\
n = 7 \\
n = 8 \\
n = 9 \\
n = 10 \\
n = 11 \\
n = 12 \\
n = 13 \\
n \geq 14 \\
\text{if } m \geq 8 \text{ and} \\
n \leq 0 \\
n \leq 1 \\
n = 2 \\
n = 3 \\
n = 4 \\
n = 5 \\
6 \leq n \leq m-2 \\
n = m-1 \\
n = m \\
n = m+1 \\
n = m+2 \\
n = m+3 \\
n = m+4 \\
n = m+5 \\
n = m+6 \\
n \geq m+7
\end{array}
\right.
\end{aligned}$$

$$m \heartsuit (-5)_n = A^{(-5)}(n) = (A^{(-4)} \circ s \circ A^{(-4)-1})(n) =$$

$n+1$	$\wedge 1$	if $m = 2$ and
$n+1$	$\equiv 2$	$n \wedge 0$
$n+2$	$\equiv 4$	$n \equiv 1$
$n+3$	$\equiv 6$	$n \equiv 2$
$n+1$	$\equiv 5$	$n \equiv 3$
$n+2$	$\equiv 7$	$n \equiv 4$
$n+3$	$\equiv 9$	$n \equiv 5$
$n+1$	$\equiv 8$	$n \equiv 6$
$n-5$	$\equiv 3$	$n \equiv 7$
$n+1$	$\vee 10$	$n \equiv 8$
		$n \vee 9$
		if $m = 3$ and
$n+1$	$\wedge 1$	$n \wedge 0$
$n+2$	$\equiv 3$	$n \equiv 1$
$n+6$	$\equiv 8$	$n \equiv 2$
$n+6$	$\equiv 9$	$n \equiv 3$
$n+2$	$\equiv 6$	$n \equiv 4$
$n+2$	$\equiv 7$	$n \equiv 5$
$n-4$	$\equiv 2$	$n \equiv 6$
$n+3$	$\equiv 10$	$n \equiv 7$
$n-3$	$\equiv 5$	$n \equiv 8$
$n-5$	$\equiv 4$	$n \equiv 9$
$n+1$	$\vee 11$	$n \vee 10$
		if $m = 4$ and
$n+1$	$\wedge 1$	$n \wedge 0$
$n+3$	$\equiv 4$	$n \equiv 1$
$n+7$	$\equiv 9$	$n \equiv 2$
$n+8$	$\equiv 11$	$n \equiv 3$
$n+6$	$\equiv 10$	$n \equiv 4$
$n+3$	$\equiv 8$	$n \equiv 5$
$n-1$	$\equiv 5$	$n \equiv 6$
$n-4$	$\equiv 3$	$n \equiv 7$
$n-6$	$\equiv 2$	$n \equiv 8$
$n-2$	$\equiv 7$	$n \equiv 9$
$n-4$	$\equiv 6$	$n \equiv 10$
$n+1$	$\vee 12$	$n \vee 11$
		if $m = 5$ and
$n+1$	$\wedge 1$	$n \wedge 0$
$n+4$	$\equiv 5$	$n \equiv 1$
$n+4$	$\equiv 6$	$n \equiv 2$
$n+9$	$\equiv 12$	$n \equiv 3$
$n-2$	$\equiv 2$	$n \equiv 4$
$n+6$	$\equiv 11$	$n \equiv 5$
$n+3$	$\equiv 9$	$n \equiv 6$
$n-4$	$\equiv 3$	$n \equiv 7$
$n+2$	$\equiv 10$	$n \equiv 8$
$n-1$	$\equiv 8$	$n \equiv 9$
$n-3$	$\equiv 7$	$n \equiv 10$
$n-7$	$\equiv 4$	$n \equiv 11$
$n+1$	$\vee 13$	$n \vee 12$
		if $m = 6$ and
$n+1$	$\wedge 1$	$n \wedge 0$
$n+5$	$\equiv 6$	$n \equiv 1$
$n+6$	$\equiv 8$	$n \equiv 2$
$n+10$	$\equiv 13$	$n \equiv 3$
$n-2$	$\equiv 2$	$n \equiv 4$
$n-2$	$\equiv 3$	$n \equiv 5$
$n+6$	$\equiv 12$	$n \equiv 6$
$n+3$	$\equiv 10$	$n \equiv 7$
$n-1$	$\equiv 7$	$n \equiv 8$
$n+2$	$\equiv 11$	$n \equiv 9$
$n-1$	$\equiv 9$	$n \equiv 10$
$n-6$	$\equiv 5$	$n \equiv 11$
$n-8$	$\equiv 4$	$n \equiv 12$
$n+1$	$\vee 14$	$n \vee 13$
		if $m = 7$ and
$n+1$	$\wedge 1$	$n \wedge 0$
$n+6$	$\equiv 7$	$n \equiv 1$
$n+4$	$\equiv 6$	$n \equiv 2$
$n+11$	$\equiv 14$	$n \equiv 3$
$n-2$	$\equiv 2$	$n \equiv 4$
$n-2$	$\equiv 3$	$n \equiv 5$
$n+3$	$\equiv 9$	$n \equiv 6$
$n+6$	$\equiv 13$	$n \equiv 7$
$n+3$	$\equiv 11$	$n \equiv 8$
$n-1$	$\equiv 8$	$n \equiv 9$
$n+2$	$\equiv 12$	$n \equiv 10$
$n-1$	$\equiv 10$	$n \equiv 11$
$n-7$	$\equiv 5$	$n \equiv 12$
$n-9$	$\equiv 4$	$n \equiv 13$
$n+1$	$\vee 15$	$n \vee 14$
		if $m \geq 8$ and
$n+1$	$\wedge 1$	$n \wedge 0$
$m$	$\equiv m$	$n \equiv 1$
$n+4$	$\equiv 6$	$n \equiv 2$
$m+7$	$\equiv m+7$	$n \equiv 3$
$n-2$	$\equiv 2$	$n \equiv 4$
$n-2$	$\equiv 3$	$n \equiv 5$
$n+1$	$\vee 7, \wedge m-1$	$6 \wedge n \leq m-2$
$n+3$	$\equiv m+2$	$n \equiv m-1$
$n+6$	$\equiv m+6$	$n \equiv m$
$n+3$	$\equiv m+4$	$n \equiv m+1$
$n-1$	$\equiv m+1$	$n \equiv m+2$
$n+2$	$\equiv m+5$	$n \equiv m+3$
$n-1$	$\equiv m+3$	$n \equiv m+4$
5	$\equiv 5$	$n \equiv m+5$
4	$\equiv 4$	$n \equiv m+6$
$n+1$	$\vee m+8$	$n \vee m+7$

## Postscript

I am translating ”ハイパー演算の整数拡張” which I wrote in Japanese last year.

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